# ON GÖDEL'S SECOND INCOMPLETENESS THEOREM 

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Abstract. A short proof of Gödel's Second Incompleteness Theorem

Gödel's Second Incompleteness Theorem states that no sufficiently strong consistent mathematical theory can prove its own consistency ([1]). In this note we give a short proof of the theorem.

Theorem. It is unprovable in set theory (unless it is inconsistent) that there exists a model of set theory.

Proof. Assume that set theory is consistent and that it proves that a model of set theory exists. Let $\Sigma$ be a finite set of axioms sufficiently strong to formulate the concepts "model" and "satisfies" $\Gamma$ and to prove the existence of a model of set theory. For the rest of the proofTa model means a model of $\Sigma \Gamma$ and letters $M, N$ denote models. If $m$ is a set with a binary relation $\Gamma \epsilon^{m}$ denotes that relation. If $N \models(E$ is a relation $) \Gamma$ then $E^{*}$ is the relation consisting of all pairs $(x, y)$ such that $N \models x E y$.

If $M$ and $N$ are models we define

$$
M<N \text { if there exists some } m \in N \text { such that } \epsilon^{M}=\left(\epsilon^{m}\right)^{*} .
$$

Informally $\Gamma M<N$ means that $M$ is $\Gamma$ in the real world $\Gamma$ the structure that $N$ thinks $m$ is. If this is the case then for every sentence $\sigma$,

$$
\begin{equation*}
M \models \sigma \quad \text { if and only if } \quad N \models(m \models \sigma) . \tag{1}
\end{equation*}
$$

[^0]In particular $\Gamma N \neq(m$ is a model $)$. Also $\Gamma$ if $N \models$ ( $m$ is a model ), then $\left(\epsilon^{m}\right)^{*}$ is the $\in-$ relation of some model $M<N$. It follows that

$$
\begin{equation*}
\text { if } M_{1}<M_{2} \text { and } M_{2}<M_{3} \text { then } M_{1}<M_{3} . \tag{2}
\end{equation*}
$$

Let us consider some fixed coding of formulas by numbers (Gödel numbering) Гand let $S_{n}$ be the name for the $n$th definable set of numbers.

Definition. $S$ is the set of all numbers $n$ with the property that there is a model $M$ such that $M \models n \notin S_{n}$.

Let $k$ be the Gödel number of $S$ and let $A$ be the sentence " $k \in S$ ". Then the following equivalence is provable in $\Sigma$ :

$$
\begin{equation*}
A \leftrightarrow \exists M(M \models \neg A) \tag{3}
\end{equation*}
$$

By (1) $\Gamma$ if $M$ is any model then

$$
\begin{equation*}
M \models A \leftrightarrow \exists N<M(N \models \neg A) \tag{4}
\end{equation*}
$$

We say that $M$ is positive if $M \models A \Gamma$ and negative otherwise. As a consequence of (4) $\Gamma$ if $M$ is negative then all $N<M$ are positive.

Since $\Sigma$ is consistent and proves that a model exists $\Gamma$ we have
there exists a model
and also (using (1))
for every model $M$ there exists a model $N<M$.

Toward a contradiction $\Gamma$ let $M_{1}$ be a model $\Gamma$ by (5). If $M_{1}$ is positive $\Gamma$ there is $\Gamma$ by (4) $\Gamma$ a negative model $M_{2}<M_{1}$; otherwise let $M_{2}=M_{1}$. By (6) there is a model $M_{3}<M_{2}$ and since $M_{2}$ is negative $\Gamma M_{3}$ is positive. By (4) there is a negative $M_{4}<M_{3}$ and we have $M_{4}<M_{2}$ by (2). A contradiction.

Remark 1. The sentence $A$ in (3) is the analog of Gödel's "I am unprovable". Another way to obtain $A$ is as follows: A property is a formula of the language of set theory with
one free variable. Let $p$ be the property (of properties $q$ ) $\exists M M \models \neg q(q)$, and let $A=p(p)$. Then (3) holds.

Remark 2. Even though our proof of Gödel's Theorem uses the Completeness TheoremГit can be modified to apply to weaker theories such as Peano Arithmetic (PA). To prove that PA does not prove its own consistency (unless it is inconsistent) $\Gamma$ we argue as follows:

Assume that PA is consistent and that "PA is consistent" is provable in PA. There is a conservative extension $\Gamma$ of PA in which the Completeness Theorem is provable ([2] $\Gamma$ Theorem $5.5 \Gamma$ page 443 ) and moreover $\Gamma$ PA $\vdash(\Gamma$ is a conservative extension of PA). Therefore $\Gamma \vdash$ ( $\Gamma$ is a conservative extension of a consistent theory) and thus proves its own consistency. Consequently $\Gamma \Gamma$ proves that $\Gamma$ has a model.

Now let $\Sigma$ be a sufficiently strong finite subset of $\Gamma$ that proves that $\Sigma$ has a model; the proof above leads to a contradiction.

## References

1. Kurt Gödel, Über formal unentscheidbare Sätze der Principia mathematica und verwandter Systeme I, Monatsh. Math. Phys. 38 (1931), 173-198.
2. Gaisi Takeuti, Proof Theory, Second Edition, Appendix by S. Simpson, North-Holland Publ. Co., Amsterdam, 1987.

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[^0]:    Supported in part by NSF grant DMS-8918299. I am grateful to Andreas Blass for comments and improvements

