

Theorems of Church and Trakhtenbrot

Stephen G. Simpson

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1 Describing a Run of a Program

Let P be a register machine program. Let R_1, \dots, R_k be the registers that are used in P . We imagine the registers as boxes which hold a finite number of marbles. Let I_1, \dots, I_l be the instructions of P . Each of I_1, \dots, I_l is either an increment instruction (add one marble) or a decrement instruction (branch on empty, remove one marble). We assume that P starts by executing instruction I_1 at time zero. We assume that time is discrete and linearly ordered. Let I_0 be a stop instruction. We assume that time stops if and when I_0 is reached.

We are going to write a predicate calculus sentence F_P which describes a run of P . To write our sentence, we use a language with the following $k + l + 7$ predicates:

Ixy : x is identical to y

Tx : x is a time instant

Zx : x is time instant zero

Lxy : x and y are time instants and y is later than x

Nxy : x and y are time instants and y is immediately after x

Mx : x is a marble

$R_i xy$: marble x is in box R_i at time y , for $1 \leq i \leq k$

$I_m x$: instruction I_m is executed at time x , for $0 \leq m \leq l$

Let F_P be the conjunction of the following axioms:

1. identity axioms:

- (a) $\forall x Ixx$
- (b) $\forall x \forall y (Ixy \Rightarrow \forall z (Ixz \Leftrightarrow Iyz))$
- (c) $\forall x \forall y (Ixy \Rightarrow \forall z (Izx \Leftrightarrow Izy))$
- (d) $\forall x \forall y (Ixy \Rightarrow (Tx \Leftrightarrow Ty))$
- (e) $\forall x \forall y (Ixy \Rightarrow (Zx \Leftrightarrow Zy))$
- (f) $\forall x \forall y (Ixy \Rightarrow \forall z (Nxz \Leftrightarrow Nyz))$
- (g) $\forall x \forall y (Ixy \Rightarrow \forall z (Nzx \Leftrightarrow Nzy))$
- (h) $\forall x \forall y (Ixy \Rightarrow \forall z (Lxz \Leftrightarrow Lyz))$
- (i) $\forall x \forall y (Ixy \Rightarrow \forall z (Lzx \Leftrightarrow Lzy))$
- (j) $\forall x \forall y (Ixy \Rightarrow (Mx \Leftrightarrow My))$
- (k) $\forall x \forall y (Ixy \Rightarrow \forall z (R_ixz \Leftrightarrow R_izy))$, for $1 \leq i \leq k$
- (l) $\forall x \forall y (Ixy \Rightarrow \forall z (R_izx \Leftrightarrow R_izy))$, for $1 \leq i \leq k$
- (m) $\forall x \forall y (Ixy \Rightarrow (I_mx \Leftrightarrow I_my))$, for $0 \leq m \leq l$

2. structural axioms:

- (a) $\forall x (Mx \vee Tx)$
- (b) $\neg \exists x (Mx \wedge Tx)$
- (c) $\forall x \forall u (R_ixu \Rightarrow (Mu \wedge Tx))$, for $1 \leq i \leq k$
- (d) $\neg \exists x \exists u (R_ixu \wedge R_jxu)$, for $1 \leq i < j \leq k$
- (e) $\forall x \forall y ((Tx \wedge Ty \wedge \neg Ixy) \Leftrightarrow (Lxy \vee Lyx))$
- (f) $\forall x \forall y \forall z ((Lxy \wedge Lyz) \Rightarrow Lxz)$
- (g) $\forall x \forall y (Nxy \Leftrightarrow (Lxy \wedge \neg \exists z (Lxz \wedge Lzy)))$
- (h) $\forall x (Tx \Leftrightarrow (I_0x \vee I_1x \vee \dots \vee I_lx))$
- (i) $\neg \exists x (I_mx \wedge I_nx)$, for $0 \leq m < n \leq l$
- (j) $\exists x Zx$
- (k) $\forall x (Zx \Rightarrow (I_1x \wedge \neg \exists y Lyx))$

- (l) $\forall x (I_m x \Rightarrow \exists y Nxy)$, for $1 \leq m \leq l$
- (m) $\forall x (I_0 x \Rightarrow \neg \exists y Lxy)$

3. axioms describing increment instructions:

For each m in the range $1 \leq m \leq l$, if I_m is an increment instruction then it is of the form

increment register R_i and go to instruction I_n

where $1 \leq i \leq k$ and $0 \leq n \leq l$, and we have:

- (a) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow I_n y)$
- (b) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u (R_i u x \Rightarrow R_i u y))$
- (c) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \exists u (R_i u y \wedge \neg R_i u x))$
- (d) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u \forall v ((R_i u y \wedge \neg R_i u x \wedge R_i v y \wedge \neg R_i v x) \Rightarrow Iuv))$
- (e) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u (R_j u x \Leftrightarrow R_j u y))$, for $1 \leq j \leq k$, $j \neq i$

4. axioms describing decrement instructions:

For each m in the range $1 \leq m \leq l$, if I_m is a decrement instruction then it is of the form

if R_i empty go to I_{n_0} , otherwise decrement R_i and go to I_{n_1}

where $1 \leq i \leq k$ and $0 \leq n_0 \leq l$ and $0 \leq n_1 \leq l$, and we have:

- (a) $\forall x \forall y ((I_m x \wedge Nxy \wedge \neg \exists u R_i u x) \Rightarrow I_{n_0} y)$
- (b) $\forall x \forall y ((I_m x \wedge Nxy \wedge \exists u R_i u x) \Rightarrow I_{n_1} y)$
- (c) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u (R_i u y \Rightarrow R_i u x))$
- (d) $\forall x \forall y ((I_m x \wedge Nxy \wedge \exists u R_i u x) \Rightarrow \exists u (R_i u x \wedge \neg R_i u y))$
- (e) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u \forall v ((R_i u x \wedge \neg R_i u y \wedge R_i v x \wedge \neg R_i v y) \Rightarrow Iuv))$
- (f) $\forall x \forall y ((I_m x \wedge Nxy) \Rightarrow \forall u (R_j u x \Leftrightarrow R_j u y))$, for $1 \leq j \leq k$, $j \neq i$

Note that F_P describes a run of the program P , but we have not specified the initial contents of the registers.

2 Unsolvability

The *Halting Problem* is the problem of deciding whether a given register machine program started with all registers empty eventually stops. It is well known that the Halting Problem is unsolvable.

The *Validity Problem* is the problem of deciding whether a given predicate calculus sentence is valid. The *Satisfiability Problem* is the problem of deciding whether a given predicate calculus sentence is satisfiable. We are going to show that the Validity Problem and the Satisfiability Problem are unsolvable. This will be accomplished by reducing the Halting Problem to them. In other words, we shall show that if either of them were solvable, then the Halting Problem would be solvable.

Given a register machine program P , let $P(0)$ be the unique run of P starting with all registers empty. Let F_P be as in Section 1, and let $F_P(0)$ be F_P conjuncted with

$$\forall x (Zx \Rightarrow \neg \exists u (R_1ux \vee \cdots \vee R_kux)) .$$

Thus $F_P(0)$ describes $P(0)$.

Let S be the sentence $\exists x I_0x$. If $P(0)$ eventually stops, then clearly $F_P(0)$ is satisfiable in a finite domain, and any domain satisfying $F_P(0)$ is necessarily finite and satisfies S . Hence in this case $F_P(0) \Rightarrow S$ is valid.

On the other hand, if $P(0)$ does not eventually stop, then $F_P(0)$ is not satisfiable in any finite domain, but $F_P(0) \wedge \neg S$ is satisfiable in an infinite domain. Thus in this case $F_P(0) \Rightarrow S$ is not valid.

We have proved:

Theorem 1 (Church's Theorem) *The Satisfiability Problem and the Validity Problem are unsolvable.*

Theorem 2 (Trakhtenbrot's Theorem) *The problem of satisfiability in a finite domain is unsolvable. The problem of validity in finite domains is unsolvable.*

3 Recursive Inseparability

Let V be the set of Gödel numbers of valid sentences, and let V_{fin} be the set of Gödel numbers of sentences which are valid in all finite domains. Note

that $V \subseteq V_{\text{fin}}$. Theorems 1 and 2 can be rephrased by saying that neither V nor V_{fin} is recursive.

We shall now prove the following stronger result, also due to Trakhtenbrot.

Theorem 3 *There is no recursive set X such that $V \subseteq X \subseteq V_{\text{fin}}$.*

Remark. By the Gödel Completeness Theorem, V is recursively enumerable, *i.e.*, Σ_1^0 . It can also be shown that V_{fin} is co-recursively enumerable, *i.e.*, Π_1^0 (this is straightforward). Thus Theorem 3 implies that V and the complement of V_{fin} form a recursively inseparable pair of recursively enumerable sets.

In general, a pair of recursively enumerable sets A and B is said to be *recursively inseparable* if $A \cap B = \emptyset$ and there is no recursive set X such that $A \subseteq X$ and $X \cap B = \emptyset$. The existence of a recursively inseparable pair of recursively enumerable sets is easily proved by a diagonal argument. For example, we may take $A = \{e \mid \varphi_e^{(1)}(e) \simeq 0\}$ and $B = \{e \mid \varphi_e^{(1)}(e) \simeq 1\}$. If X were a recursive set separating A from B , then letting e be an index of the characteristic function of X we would have $e \in X$ if and only if $e \notin X$, a contradiction.

In order to prove Theorem 3, we shall slightly modify the construction of Section 2.

Let A and B be a recursively inseparable pair of recursively enumerable sets. Let f be the partial recursive function defined by

$$f(n) \simeq \begin{cases} 0 & \text{if } n \in A, \\ 1 & \text{if } n \in B, \\ \text{undefined} & \text{otherwise.} \end{cases}$$

Let P be a register machine program which computes f . Let $P(n)$ be the unique run of P starting with n in R_1 and all other registers empty. Let F_P be as before, and let $F_P(n)$ be F_P conjuncted with

$$\forall x (Zx \Rightarrow \exists \text{ exactly } n \text{ } u \text{ such that } R_1ux)$$

conjuncted with

$$\forall x (Zx \Rightarrow \neg \exists u (R_2ux \vee \cdots \vee R_kux)) .$$

Thus $F_P(n)$ describes $P(n)$.

Let S_0 be the sentence $\exists x (I_0x \wedge \neg \exists u R_2ux)$. If $n \in A$ then $P(n)$ eventually stops with 0 in R_2 , hence $F_P(n) \Rightarrow S_0$ is valid, hence the Gödel number of $F_P(n) \Rightarrow S_0$ belongs to V . If $n \in B$, then $P(n)$ eventually stops with 1 in R_2 , hence $F_P(n) \wedge \neg S_0$ is satisfiable in a finite domain, hence the Gödel number of $F_P(n) \Rightarrow S_0$ does not belong to V_{fin} .

We can now complete the proof of Theorem 3. If there were a recursive set X such that $V \subseteq X \subseteq V_{\text{fin}}$, then

$$\{n : \text{the Gödel number of } F_P(n) \Rightarrow S_0 \text{ belongs to } X\}$$

would be a recursive set which separates A from B . Since A and B are recursively inseparable, Theorem 3 follows.